Splitting algorithms for non-smooth convex optimization<sup>1</sup>: Review and application to Mean Field  $Games^2$ 

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Split/Unsplit





#### Mean Field Games $(MFG)^3$ : Static example

- N players choose their positions on a set Q (compact).
- $\mathcal{P}(Q)$  is the set of Borel probability measures.
- They minimize their distance to a place  $P \in Q$ .
- Players are congestion-averse.
- The cost of player i can be modeled by

$$f_i(x_1, \dots, x_i, \dots, x_N) = \alpha |x_i - P| - \frac{\beta}{N-1} \sum_{j \neq i} |x_j - x_i|$$
$$= \alpha |x_i - P| - \beta \int_Q |x - x_i| d\left(\frac{1}{N-1} \sum_{j \neq i} \delta_{x_j}\right)$$
$$= f\left(x_i, \frac{1}{N-1} \sum_{j \neq i} \delta_{x_j}\right).$$

<sup>3</sup>J.-M. Lasry, P.-L. Lions. Mean field games. Jpn. J. Math. 2007

M. Huang, R. P. Malhamé, P. E. Caines. Large population stochastic dynamic games. Commun. Inf. Syst. 2006.  $\langle \Box \rangle \langle \overline{\Box} \rangle \langle \overline{\Xi} \rangle \langle \overline{\Xi} \rangle$ 

• Suppose that for each N,  $(\bar{x}_1^N, \ldots, \bar{x}_N^N)$  is a Nash equilibrium of the previous game, i.e., for every  $i \in \{1, \ldots, N\}$ ,

Splitting algorithms

$$(\forall x_i \in Q) \quad f\left(\bar{x}_i^N, \frac{1}{N-1}\sum_{j \neq i} \delta_{\bar{x}_j^N}\right) \le f\left(x_i, \frac{1}{N-1}\sum_{j \neq i} \delta_{\bar{x}_j^N}\right)$$

• Then,  $\exists \ \bar{m} \in \mathcal{P}(Q)$  such that, up to some sub-sequence,

$$\frac{1}{N}\sum_{i=1}^N \delta_{\bar{x}_i^N} \stackrel{*}{\rightharpoonup} \bar{m}.$$

• The equilibrium  $\bar{m}$  satisfies the fixed point equation

 $\operatorname{supp}(\bar{m}) \subseteq \operatorname{argmin} \{ f(x, \bar{m}) \mid x \in Q \}.$ 

Motivation

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Dynamic & deterministic case

• Differential game with N players, where Player i minimizes

Splitting algorithms

$$\begin{split} \int_0^T \left[ \frac{|\alpha(t)|^2}{2} + f\left(x_i(t), \frac{1}{N-1} \sum_{j \neq i} \delta_{x_j(t)}\right) \right] dt + g\left(x_i(T), \frac{1}{N-1} \sum_{j \neq i} \delta_{x_j(T)}\right) \\ \text{s.t.} \quad \dot{x}_i(t) = \alpha(t) \quad \forall t \in [0, T], \\ x_i(0) = \bar{x}_{0, i}^N. \end{split}$$

• Suppose that  $(\bar{x}_1^N, \dots, \bar{x}_N^N)$  is a Nash equilibrium and that

$$\frac{1}{N}\sum_{i=1}^N \delta_{\bar{x}_{0,i}^N} \stackrel{*}{\rightharpoonup} \underline{m_0}.$$

Then for each  $t \in [0, T], \exists m(t) \in \mathcal{P}(Q)$  such that

$$\frac{1}{N}\sum_{i=1}^N \delta_{\bar{x}_i^N(t)} \stackrel{*}{\rightharpoonup} m(t)$$

Numerical experiences

Motivation

Dynamic & deterministic case

• Any equilibrium m solves the MFG system

Splitting algorithms

$$(HJB) \quad -\partial_t u + \frac{|\nabla u|^2}{2} = f(x, m(t))$$
$$u(x, T) = g(x, m(T)),$$
$$(FP) \quad \partial_t m - \operatorname{div}(m\nabla u) = 0,$$
$$m(0) = m_0.$$

• At (x, t) the solution u of the HJB equation is given by

$$u(x,t) = \inf_{\alpha} \int_{t}^{T} \left[ \frac{|\alpha(s)|^2}{2} + f(x(s), \mathbf{m}(s)) \right] ds + g(x(T), \mathbf{m}(T))$$
  
s.t.  $\dot{x}(s) = \alpha(s) \quad \forall s \in (t,T),$   
 $x(t) = x.$ 

Numerical experiences

Motivation

Dynamic & stochastic case

$$\int_0^T \left[ \frac{|\alpha(t)|^2}{2} + f\left(x_i(t), \frac{1}{N-1} \sum_{j \neq i} \delta_{x_j(t)}\right) \right] dt + g\left(x_i(T), \frac{1}{N-1} \sum_{j \neq i} \delta_{x_j(T)}\right)$$
  
s.t.  $dx_i(t) = \alpha(t) dt + \sigma dW_i(t) \quad \forall t \in [0, T],$   
 $x_i(0) = \bar{x}_{0,i}^N.$ 

As before,  $\frac{1}{N}\sum_{i=1}^{N} \delta_{\bar{x}_{i}^{N}(t)} \to m(t)$  which now solves  $(t \in [0, T], x \in Q)$ :

MFG

$$-\partial_t u - \frac{\sigma^2}{2} \Delta u + \frac{|\nabla u|^2}{2} = f(x, m(t))$$
$$u(x, T) = g(x, m(T))$$
$$\partial_t m - \frac{\sigma^2}{2} \Delta m - \operatorname{div}(m \nabla u) = 0$$
$$m(0) = m_0$$

## Stationary MFG

Motivation

• The ergodic counterpart of MFG is  $(x \in Q)$ 

Splitting algorithms

Stationary MFG (SMFG)

$$\begin{aligned} &-\frac{\sigma^2}{2}\Delta u + \frac{|\nabla u|^2}{2} + \lambda = f(x,m), \\ &-\frac{\sigma^2}{2}\Delta m - \operatorname{div}(m\nabla u) = 0, \\ &\int_Q u(x)dx = 0, \quad m \ge 0, \quad \int_Q m(x)dx = 1. \end{aligned}$$

Numerical experiences

• The solution  $(u, m, \lambda)$  of SMFG describes the long time average of solutions  $(u^T, m^T)$  of MFG as  $T \to \infty^4$ .

<sup>&</sup>lt;sup>4</sup>P. Cardaliaguet, J.-M. Lasry, P.-L. Lions, A. Porretta, Long time average of mean field games with a nonlocal coupling, SIAM J. Control Optim., 2013 (□) (□) (□) (□)

Stationary MFG with local couplings

• In this talk we focus on local interactions f(x,m) = f(x,m(x)).

Splitting algorithms

• We set 
$$Q = \mathbb{T}^2$$
 and  $\nu = \sigma^2/2$ .

• We study numerical methods for solving

#### SMFG

Motivation

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$$\begin{aligned} -\nu\Delta u(x) &+ \frac{1}{2} |\nabla u(x)|^2 + \lambda = f(x, m(x)) \quad \text{in } \mathbb{T}^2 \\ &-\nu\Delta m(x) - \operatorname{div} \left( m(x) \nabla u(x) \right) = 0 \quad \text{in } \mathbb{T}^2 \\ m &\ge 0, \quad \int_{\mathbb{T}^2} m(x) dx = 1, \quad \int_{\mathbb{T}^2} u(x) dx = 0. \end{aligned}$$

Motivation

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Numerical experiences

## Discrete SMFG

DSMFG (Achdou & Capuzzo Dolcetta, 2010)

$$-\nu(\Delta_h u^h)_{i,j} + \frac{1}{2} |[\widehat{D_h u^h}]_{i,j}|^2 + \lambda^h = f(x_{i,j}, m_{i,j}^h) \quad \forall 0 \le i,j \le N_h - 1$$
$$-\nu(\Delta_h m^h)_{i,j} - \left(\operatorname{div}_h(m^h[\widehat{D_h u^h}])\right)_{i,j} = 0 \quad \forall 0 \le i,j \le N_h - 1$$
$$m_{i,j}^h \ge 0, \ h^2 \sum_{i,j} m_{i,j}^h = 1, \ \sum_{i,j} u_{i,j}^h = 0.$$

• 
$$h > 0, N_h = 1/h, \mathcal{M}_h = \mathbb{R}^{N_h \times N_h}, \mathcal{W}_h = \mathbb{R}^{4(N_h \times N_h)}$$

•  $\widehat{[D_h u]}_{i,j} = ((D_1 u)_{i,j}^-, (D_1 u)_{i-1,j}^+, (D_2 u)_{i,j}^-, (-D_2 u)_{i,j-1}^+) \in \mathbb{R}^4$ , where  $(D_1 u)_{i,j} := \frac{u_{i+1,j} - u_{i,j}}{h}, (D_2 u)_{i,j} := \frac{u_{i,j+1} - u_{i,j}}{h}.$ 

•  $\Delta_h$  and div<sub>h</sub> are linear operators defined by  $(\Delta_h m)_{i,j} := -\frac{1}{h^2} (4m_{i,j} - m_{i+1,j} - m_{i-1,j} - m_{i,j+1} - m_{i,j-1})$  $(\operatorname{div}_h(w))_{i,j} := (D_1 w^1)_{i-1,j} + (D_1 w^2)_{i,j} + (D_2 w^3)_{i,j-1} + (D_2 w^4)_{i,j}.$  SMFG discretization

Motivation

• If  $\nu > 0$ ,  $f(x, \cdot)$  is increasing and we suppose that the stationary system admits a unique classical solution, in Achdou, Camilli & Capuzzo Dolcetta (2013) the convergence of DSMFG (unif- $L^2$ ) to the unique solution to the stationary system as  $h \to 0$  is proved.

Splitting algorithms

- To solve the discretized system, Newton's method can be used (Achdou & Capuzzo Dolcetta, 2010; Achdou & Perez, 2012; Cacace & Camilli, 2016) if the initial guess is close enough to the solution.
- The performance of Newton's method depends heavily on the values of  $\nu$ : for small values of  $\nu$  the convergence is much slower and cannot be guaranteed in general since  $m^h$ can become negative.

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## Variational Approach

$$b(m,w) = \begin{cases} \frac{|w|^2}{2m}, & \text{if } m > 0;\\ 0, & \text{if } (m,w) = (0,0); \\ +\infty, & \text{otherwise}, \end{cases} \quad F(x,m) = \begin{cases} \int_0^m f(x,m')dm', & \text{if } m \ge 0;\\ +\infty, & \text{otherwise}. \end{cases}$$

The SMFG is (formally) the FOC of the optimization problem

#### Optimization Problem (P)

$$\begin{split} &\inf_{m,w}\int_{\mathbb{T}^2}\left[b(m(x),w(x))+F(x,m(x))\right]dx\\ \text{s.t}\quad \begin{cases} -\nu\Delta m(x)+\operatorname{div}\bigl(w(x)\bigr)=0, \quad \text{in } \mathbb{T}^2\\ \int_{\mathbb{T}^2}m(x)dx=1, \end{cases} \end{split}$$

where u and  $\lambda$  are Lagrange multipliers and  $w = -m\nabla u$  (see Lasry & Lions, 2007).

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Motivation

## Goal of this talk...

- Provide variational formulation  $(P_h)$  of DSMFG (Achdou & Capuzzo Dolcetta, 2010) in order to propose numerical approximations of the SMFG.
- Review of proximal splitting methods for solving  $(P_h)$ .
  - Benamou & Carlier (2015) and Benamou, Carlier & Santambrogio (2017): FE discretization of the dynamic MFG via Augmented Lagrangian method (ADMM) case  $\nu = 0$ .
  - And reev (2017): ADMM with preconditioners for  $\nu>0$  in dynamic case. Another discretization.
  - Papadakis, Peyré & Oudet (2014): optimal transport (F = 0)and  $\nu = 0$  with centered grid.
- Connect first-order methods and fixed point iterations.
- Propose a new **projected** proximal splitting method for solving  $(P_h)$ .
- Numerical experiments.

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**3** Splitting algorithms: review







$$\begin{array}{c|c} \hline \text{Motivation} & (P_h) & \text{Splitting algorithms} & \text{Split/Unsplit} & \text{Numerical experienc} \\ \hline \text{Optimization Problem } (P_h) \\ \hline \text{Optimization Problem } (P_h) \\ \hline \begin{array}{c} \text{Discrete optimization problem } (P_h) \\ & \inf_{(m,w) \in \mathcal{M}_h \times \mathcal{W}_h} h^2 \sum_{i,j=0}^{N_h - 1} \left[ \hat{b}(m_{i,j}, w_{i,j}) + F(x_{i,j}, m_{i,j}) \right] \\ & \text{s.t.} & \begin{cases} -\nu(\Delta_h m)_{i,j} + (\operatorname{div}_h w)_{i,j} = 0, & \forall 0 \leq i, j, \leq N_h - 1 \\ h^2 \sum_{i,j} m_{i,j} = 1. \end{cases} \end{array} \right.$$

−νΔ<sub>h</sub>: M<sub>h</sub> → M<sub>h</sub> and div<sub>h</sub>: W<sub>h</sub> → M<sub>h</sub> are linear. *b*: ℝ × ℝ<sup>4</sup> is given by

$$\hat{b} \colon (m,w) \mapsto \begin{cases} \frac{|w|^2}{2m}, & \text{if } m > 0, \ w \in K, \\ 0, & \text{if } (m,w) = (0,0), \\ +\infty, & \text{otherwise.} \end{cases}$$

•  $K := \mathbb{R}_+ \times \mathbb{R}_- \times \mathbb{R}_+ \times \mathbb{R}_-$ .

Split/Unsplit

## Existence of solutions to $(P_h)$

#### Qualification condition

 $(P_h)$ 

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There exists a feasible  $(\tilde{m}, \tilde{w}) \in \mathcal{M}_h \times \mathcal{W}_h$  such that  $\tilde{w} \in \operatorname{int}(K), \quad \tilde{m}_{i,j} > 0 \ \forall (i, j).$ 

#### Existence (BA-Kalise-Silva, 2018)

For any  $\nu \geq 0$  we have

- **(** $P_h$ ) admits at least one solution and the optimal costs are finite.
- 2 Let  $(m^h, w^h)$  be a solution to  $(P_h)$ . Then, there exists  $(u^h, \lambda^h) \in \mathcal{M}_h \times \mathbb{R}$  s.t. DSMFG holds  $(w^h = m^h[\widehat{D_h u^h}])$ .

*Proof:* Qualification condition implies the existence of Lagrange multipliers. DSMFG follows from FOC.

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## $(P_h)$ 's structure

Motivation

• Assume  $f(x, \cdot)$  increasing  $(F(x, \cdot) \text{ convex})$ .

Splitting algorithms

- $\varphi: (m, w) \mapsto \sum_{i,j} \phi_{i,j}(m_{i,j}, w_{i,j})$  where,  $\forall 0 \le i, j \le N_h 1$ ,  $\phi_{i,j}(m_{i,j}, w_{i,j}) = \hat{b}(m_{i,j}, w_{i,j}) + F(x_{i,j}, m_{i,j})$  is proper, convex, l.s.c., non-smooth.
- Denote  $-\nu\Delta_h = A$  and  $\operatorname{div}_h = B$ .

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#### Reformulation of $(P_h)$

 $(P_h)$ 

$$\min_{\substack{(m,w)\in\mathcal{M}_h\times\mathcal{W}_h}}\varphi(m,w)$$
  
i.t.  $\Xi(m,w)=(0,1),$ 

where

$$\Xi = \begin{bmatrix} A & B \\ h^2 \mathbf{1}^\top & 0 \end{bmatrix}$$

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#### Reformulation of $(P_h)$

$$\min_{\substack{(m,w)\in\mathcal{M}_h\times\mathcal{W}_h}}\varphi(m,w)$$
  
s.t.  $\Xi(m,w)=(0,1).$ 

•  $(P_h)$  is equivalent to

$$(P_h) \quad \min_{(m,w)\in\mathcal{M}_h\times\mathcal{W}_h} \Phi(m,w) = \varphi(m,w) + \psi(L(m,w))$$

(Split) 
$$\begin{cases} \psi = \iota_{\{(0,1)\}} \\ L = \Xi \end{cases} \quad \text{or} \quad (\text{Unsplit}) \begin{cases} \psi = \iota_{\Xi^{-1}(0,1)} \\ L = \text{Id.} \end{cases}$$

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Split/Unsplit

Numerical experiences





**3** Splitting algorithms: review







## General convex optimization problem

#### Problem (P)

 $\underset{x \in \mathbb{R}^N}{\text{minimize}} \ \Phi(x),$ 

assuming that solution set  $Z \neq \emptyset$ .

- Φ: ℝ<sup>N</sup> → ℝ := ℝ ∪ {+∞} is
  convex : (∀x, y ∈ ℝ<sup>N</sup>)(∀λ ∈ [0, 1]) Φ(x + λ(y - x)) ≤ Φ(x) + λ(Φ(y) - Φ(x)).
  lower semicontinuous (l.s.c): (∀x ∈ ℝ<sup>N</sup>) Φ(x) ≤ lim inf<sub>y→x</sub> Φ(y).
  proper: Φ is not always +∞ and never -∞.
- Functions satisfying above conditions constitute the class  $\Gamma_0(\mathbb{R}^N)$ .

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•  $\Phi$  is convex differentiable and  $\nabla \Phi$  is  $\chi$ -Lipschitz.  $\|\nabla \Phi(x) - \nabla \Phi(y)\| \leq \chi \|x - y\|.$ 

#### Optimality conditions

$$x^*$$
 is a solution to (P)  $\Leftrightarrow 0 = \nabla \Phi(x^*).$ 

- Equivalent to, for every  $\gamma > 0$ ,  $x^* = x^* \gamma \nabla \Phi(x^*)$ .
- Equivalent to, for every  $\gamma > 0, x^*$  is a fixed point of

Gradient operator $G_{\gamma\Phi}x = x - \gamma \nabla \Phi(x).$ 

Motivation

 We can approximate solutions to (P) via fixed points iterations of the operator G<sub>γΦ</sub>, i.e.,

Splitting algorithms

Gradient method

$$x_0 \in \mathbb{R}^N$$
,  $x_{n+1} = G_{\gamma \Phi} x_n = x_n - \gamma \nabla \Phi(x_n)$ .

• When the operator  $G_{\gamma\Phi}$  is good enough for obtaining convergence ?

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Motivation

 We can approximate solutions to (P) via fixed points iterations of the operator G<sub>γΦ</sub>, i.e.,

Splitting algorithms

Gradient method

$$x_0 \in \mathbb{R}^N$$
,  $x_{n+1} = G_{\gamma \Phi} x_n = x_n - \gamma \nabla \Phi(x_n)$ .

- When the operator  $G_{\gamma\Phi}$  is good enough for obtaining convergence ?
  - $\Phi \beta$ -strongly convex  $(\Phi \beta \| \cdot \|^2/2$  convex),  $0 < \gamma < 2/\chi$ :



- $r = \max\{|1 \gamma\beta|, |1 \gamma\chi|\} < 1. \ \gamma^* = 2/(\beta + \chi)$
- Banach-Picard (linear convergence to unique solution).

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Motivation

 We can approximate solutions to (P) via fixed points iterations of the operator G<sub>γΦ</sub>, i.e.,

Splitting algorithms

Gradient method

$$x_0 \in \mathbb{R}^N$$
,  $x_{n+1} = G_{\gamma \Phi} x_n = x_n - \gamma \nabla \Phi(x_n)$ .

• When the operator  $G_{\gamma\Phi}$  is good enough for obtaining convergence ?

• 
$$0 < \gamma < 2/\chi$$
:

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Split/Unsplit

Numerical experiences

• 
$$\mu = \gamma \chi/2$$

• Convergence O(1/k) to a solution.

Motivation

 We can approximate solutions to (P) via fixed points iterations of the operator G<sub>γΦ</sub>, i.e.,

Splitting algorithms

Gradient method

$$x_0 \in \mathbb{R}^N$$
,  $x_{n+1} = G_{\gamma \Phi} x_n = x_n - \gamma \nabla \Phi(x_n)$ .

- When the operator  $G_{\gamma\Phi}$  is good enough for obtaining convergence ?
  - Variants with  $\gamma_k$  yields  $O(1/k^2)$ . (Nesterov, 1983)

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Example Grad:  $\gamma = 0.999 \in \left]0, 2/\chi\right[$ 

•  $\Phi \colon (x,y) \mapsto 0.3 \, x^2 + y^2$ 

• 
$$\nabla \Phi \colon (x,y) \mapsto (0.6 \, x, 2 \, y)$$

• 
$$\chi = 2, \, \beta = 0.6$$



## Example Grad: $\gamma = \gamma^* \in \left]0, 2/\chi\right[$

- $\Phi \colon (x,y) \mapsto 0.3 \, x^2 + y^2$
- $\nabla \Phi \colon (x,y) \mapsto (0.6 \, x, 2 \, y)$

• 
$$\chi = 2, \, \beta = 0.6, \, \gamma^* = 2/(\beta + \chi)$$





Motivation Splitting algorithms Split/Unsplit Numerical experiences  $\Phi \in \Gamma_0(\mathbb{R}^N)$  is non-smooth Problem (P)  $\underset{x \in \mathbb{R}^N}{\text{minimize}} \Phi(x)$  $\partial \Phi \colon x \mapsto \left\{ u \in \mathbb{R}^N \mid (\forall y \in \mathbb{R}^N) \ \Phi(x) + u^\top (y - x) \le \Phi(y) \right\}$ |x| $\partial |\cdot|(0) = [-1,1]$ 

Optimality conditions  $x^*$  is a solution to (P)  $\Leftrightarrow 0 \in \partial \Phi(x^*).$ 

27/51

Motivation Splitting algorithms Split/Unsplit Numerical experiences  $\Phi \in \Gamma_0(\mathbb{R}^N)$  is non-smooth Problem (P)  $\underset{x \in \mathbb{R}^N}{\text{minimize}} \Phi(x)$ **Optimality** conditions  $x^*$  is a solution to (P)  $\Leftrightarrow 0 \in \partial \Phi(x^*)$ . • Equivalent to, for every  $\gamma > 0, x^* \in x^* + \gamma \partial \Phi(x^*)$ . • Equivalent to, for every  $\gamma > 0$ ,  $x^*$  is a fixed point of

#### **Proximity operator**

$$\operatorname{prox}_{\gamma\Phi} = (\operatorname{Id} + \gamma \partial \Phi)^{-1} \colon x \mapsto \underset{y \in \mathbb{R}^N}{\operatorname{argmin}} \ \gamma \Phi(y) + \frac{1}{2} \|y - x\|^2.$$

• 
$$\operatorname{prox}_{\iota_C} = P_C$$
.

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 $\Phi \in \Gamma_0(\mathbb{R}^N)$  is non-smooth

 We can approximate solutions to (P) via fixed points iterations of the operator prox<sub>γΦ</sub>, i.e.,

Split/Unsplit

Numerical experiences

Splitting algorithms

Proximal point algorithm (PPA)

$$x_0 \in \mathbb{R}^N, \quad x_{n+1} = \operatorname{prox}_{\gamma \Phi}(x_n).$$

- When the operator  $\operatorname{prox}_{\gamma\Phi}$  is good enough ? Always (for every  $\gamma > 0$ ) (Martinet 1970-72; Rockafellar 1976).
  - $\Phi \beta$ -strongly convex:  $\operatorname{prox}_{\gamma\Phi}$  contraction  $(r = \frac{1}{1+\gamma\beta} < 1)$ .
  - In general, it is averaged nonexpansive  $(\mu = \frac{1}{2})$ . O(1/k).
  - Variations considering  $\gamma_k$  yields  $O(1/k^2)$  (Güler 1992).

Motivation

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## Example PPA: $\gamma = 1$

• 
$$\Phi \colon (x,y) \mapsto 0.3 x^2 + y^2$$

• 
$$\operatorname{prox}_{\gamma\Phi}$$
:  $(x, y) \mapsto (x/(1 + 0.6\gamma), y/(1 + 2\gamma))$ 

• 
$$\chi=2,\,\beta=0.6$$



## Example PPA: $\gamma = 100$

• 
$$\Phi \colon (x,y) \mapsto 0.3 \, x^2 + y^2$$

•  $\operatorname{prox}_{\gamma\Phi}$ :  $(x, y) \mapsto (x/(1 + 0.6\gamma), y/(1 + 2\gamma))$ 

• 
$$\chi=2,\,\beta=0.6$$





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## Our case: $\Phi = \varphi + \psi \circ L$

#### Problem (P)

$$\underset{x \in \mathbb{R}^N}{\text{minimize}} \ \Phi(x) = \varphi(x) + \psi(Lx)$$

#### • Lagrangian approach: Solve

$$\min_{Lx=y}\varphi(x) + \psi(y)$$

#### Augmented Lagrangian

$$\mathcal{L}_c(x, y, u) = \varphi(x) + \psi(y) + u^\top (Lx - y) + \frac{c}{2} \|Lx - y\|^2$$

Motivation

Our case:  $\Phi = \varphi + \psi \circ L$ 

- Under qualification conditions x solves (P) iff (x, Lx, u) is a saddle point of  $\mathcal{L}_c$ , with c > 0.
- Alternating minimization-maximization of  $\mathcal{L}_c$  we obtain:

ADMM (Glowinski-Marrocco 1975; Gabay-Mercier 1976)

$$\begin{aligned} x^{k+1} &= \operatorname{argmin}_{x} \mathcal{L}_{c}(x, y^{k}, u^{k}) \\ &= \operatorname{argmin}_{x} \left\{ \varphi(x) + u^{k^{\top}} Lx + \frac{\gamma}{2} \| Lx - y^{k} \|^{2} \right\} \\ y^{k+1} &= \operatorname{argmin}_{y} \mathcal{L}_{c}(x^{k+1}, y, u^{k}) = \operatorname{prox}_{\psi/\gamma}(u^{k}/\gamma + Lx^{k+1}) \\ u^{k+1} &= u^{k} + \gamma(Lx^{k+1} - y^{k+1}). \end{aligned}$$

- Problem: The first step is not easy to compute in general.
- Predictor-corrector proximal multiplier (PCPM) splits with additional multiplier update (Chen-Teboulle, 1994)
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Numerical experiences

## Our case: $\Phi = \varphi + \psi \circ L$

#### Problem (P)

$$\min_{x \in \mathbb{R}^N} \Phi(x) = \varphi(x) + \psi(Lx)$$

#### Fenchel-Rockafellar duality yields:

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Problem (D)

$$\min_{u \in \mathbb{R}^M} \varphi^*(-L^\top u) + \psi^*(u)$$

 $\begin{array}{c} \textbf{Fenchel conjugate} \\ \psi^* \colon u \mapsto \sup_{v \in \mathbb{R}^M} \left( u^\top v - \psi(v) \right) \end{array}$ 

•  $\psi \in \Gamma_0(\mathbb{R}^M) \Leftrightarrow \psi^* \in \Gamma_0(\mathbb{R}^M)$ •  $\partial \psi^* = (\partial \psi)^{-1}$ 

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Motivation

Our case:  $\Phi = \varphi + \psi \circ L$ 

Under qualification conditions, we have

Our case:  $\Phi = \varphi + \psi \circ L$ 

$$x \text{ solves } (P) \Leftrightarrow \begin{bmatrix} 0\\0 \end{bmatrix} \in \underbrace{\begin{bmatrix} \partial \varphi & 0\\0 & \partial \psi^* \end{bmatrix}}_{M} \underbrace{\begin{bmatrix} x\\u \end{bmatrix}}_{X} + \underbrace{\begin{bmatrix} 0 & L^\top\\-L & 0 \end{bmatrix}}_{S} \underbrace{\begin{bmatrix} x\\u \end{bmatrix}}_{X}$$
  
• for some u solving (D).

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• 
$$M(x, u) = \partial \Psi(x, u)$$
, where  $\Psi(x, u) = \varphi(x) + \psi^*(u)$ .



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• for some u solving (D).

• 
$$M(x, u) = \partial \Psi(x, u)$$
, where  $\Psi(x, u) = \varphi(x) + \psi^*(u)$ .

•  $S \neq \nabla G$ ... but it is maximally monotone:  $(SX - SY)^{\top}(X - Y) \ge 0$  and continuous.

For a monotone operator  $\boldsymbol{S}$ , we have:

$$\begin{array}{|c|c|c|c|c|}\hline \textbf{Resolvent} \\ \hline & \\ J_{\gamma \boldsymbol{S}} = (\mathrm{Id} + \gamma \boldsymbol{S})^{-1} \\ \hline & \\ G_{\gamma \boldsymbol{S}} = \mathrm{Id} - \gamma \boldsymbol{S} \\ \hline & \\ \end{array}$$

Our case:  $\Phi = \varphi + \psi \circ L$ 

Splitting approach:

• 
$$J_{\gamma M} = \operatorname{prox}_{\gamma \Psi} : (x, u) \mapsto (\operatorname{prox}_{\gamma \varphi} x, \operatorname{prox}_{\gamma \psi^*} u)$$

• 
$$G_{\gamma S} = \mathrm{Id} - \gamma S \colon (x, u) \mapsto (x - \gamma L^{\top} u, u + \gamma L x)$$

Mon.+Skew (BA-Combettes 2011)

Let  $0 < \gamma < \|L\|^{-1}, x_0 \in \mathbb{R}^N$  and  $u_0 \in \mathbb{R}^M$  and iterate

$$p_{1,n} = \operatorname{prox}_{\gamma\varphi}(x_n - \gamma L^{\top} u_n)$$
  

$$p_{2,n} = \operatorname{prox}_{\gamma\psi^*}(u_n + \gamma L x_n)$$
  

$$x_{n+1} = p_{1,n} - \gamma (L^{\top} p_{2,n} - L^{\top} u_n)$$
  

$$u_{n+1} = p_{2,n} + \gamma (L p_{1,n} - L x_n).$$

Then  $(x_n, u_n) \to (\bar{x}, \bar{u})$ , where  $\bar{x}$  solves (P) and  $\bar{u}$  solves (D).

• If  $\varphi \ \& \ \psi^*$  strongly convex: linear conv. (Tseng 2000).

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38/51

### Our case: $\Phi = \varphi + \psi \circ L$

Problem (P)

$$\min_{x \in \mathbb{R}^N} \Phi(x) = \varphi(x) + \psi(Lx)$$

Equivalent to (Condat-Vũ, 2013 & He-Yuan, 2012)

X

$$\begin{bmatrix} 0 \\ 0 \end{bmatrix} \in \underbrace{\left( \begin{bmatrix} \partial \varphi & 0 \\ 0 & \partial \psi^* \end{bmatrix} + \begin{bmatrix} 0 & L^\top \\ -L & 0 \end{bmatrix} \right)}_{\boldsymbol{A} = \boldsymbol{M} + \boldsymbol{S}} \underbrace{\begin{bmatrix} x \\ u \end{bmatrix}}_{\boldsymbol{X}}$$

•  $J_{\mathbf{A}}$  is not explicit !

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## Our case: $\Phi = \varphi + \psi \circ L$

#### Problem (P)

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- $J_{\mathbf{A}}$  is not explicit !
- Set  $U \succ 0$  symmetric and  $\langle X \mid Y \rangle_U = X^\top U Y$ .
- $\langle \boldsymbol{U}^{-1}\boldsymbol{A}\boldsymbol{X} \boldsymbol{U}^{-1}\boldsymbol{A}\boldsymbol{Y} \mid \boldsymbol{X} \boldsymbol{Y} \rangle_{\boldsymbol{U}} = (\boldsymbol{A}\boldsymbol{X} \boldsymbol{A}\boldsymbol{Y})^{\top}(\boldsymbol{X} \boldsymbol{Y}).$

•  $U^{-1}A$  is monotone under this metric: PPA.

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Motivation

Our case:  $\Phi = \varphi + \psi \circ L$ 

Specific  ${\boldsymbol U}$  allows  $J_{{\boldsymbol U}^{-1}{\boldsymbol A}}$  explicit:

Chambolle-Pock (2011)  $x_0, \bar{x}_0 \in \mathbb{R}^N \text{ and } u_0 \in \mathbb{R}^M, \ \tau, \gamma > 0 \text{ such that } \tau \gamma \|L\|^2 < 1$  $\begin{vmatrix} x^{n+1} = \operatorname{prox}_{\tau \varphi} (x^n - \tau L^\top u^n) \\ u^{n+1} = \operatorname{prox}_{\gamma \psi^*} (u^n + \gamma L(2x^{n+1} - x^n)). \end{vmatrix}$ 

Acceleration (Chambolle-Pock 2011):

- $\varphi$  and  $\psi^*$  strongly convex: linear convergence
- $\varphi$  or  $\psi^*$  strongly convex and  $\tau_k$ ,  $\gamma_k$ :  $||x_k x^*|| \le C/k$ .

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**3** Splitting algorithms: review







## $(P_h)$ 's structure

Reformulation of  $(P_h)$ 

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$$\min_{\substack{(m,w)\in\mathcal{M}_h\times\mathcal{W}_h}}\varphi(m,w)$$
  
s.t.  $\Xi(m,w)=(0,1),$ 

$$(P_h) \quad \min_{(m,w)\in\mathcal{M}_h\times\mathcal{W}_h} \Phi(m,w) = \varphi(m,w) + \psi(L(m,w))$$

(Split) 
$$\begin{cases} \psi = \iota_{\{(0,1)\}} \\ L = \Xi \end{cases} \quad \text{or} \quad (\text{Unsplit}) \begin{cases} \psi = \iota_{\Xi^{-1}(0,1)} \\ L = \text{Id.} \end{cases}$$

All previous methods implement the dual variable update

$$u^{n+1} = \operatorname{prox}_{\gamma\psi^*}(u^n + \gamma L\widetilde{x}^n).$$

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Motivation

# Split/unsplit approaches

### "Split" decomposition:

- $\psi = \iota_{\{(0,1)\}}$ ,  $\operatorname{prox}_{\gamma\psi^*} = \operatorname{Id} \gamma(0,1)$ , and  $L = \Xi$ .
- Then, all previous methods include a Lagrange multiplier step of the form

$$u^{k+1} = u^k + \gamma(\Xi x^k - (0, 1))$$

- The primal iterates  $(x_k)_{k\in\mathbb{N}}$  are not feasible !
- Very slow...

#### "Unsplit" decomposition:

- $\psi = \iota_{\{\Xi^{-1}(0,1)\}}$  and for computing  $\operatorname{prox}_{\gamma\psi^*}$  or  $\operatorname{prox}_{\gamma\psi}$  we need to invert  $\Xi\Xi^*$  (more precisely  $\nu^2 A A^* + B B^*$ ).
- Depending on the parameter  $\nu$ , this matrix can be very bad conditioned and difficult to invert.

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#### Numerical experiences

## Projected Chambolle-Pock splitting<sup>5</sup>

We avoid matrix inversions along with ensuring primal iterates to satisfy some of the constraints.

Projected Chambolle-Pock (PCP)

Let  $x_0 \in \mathcal{H}$ ,  $u_0 \in \mathcal{G}$  and  $\sigma \tau ||L||^2 < 1$ .

$$(\forall n \in \mathbb{N}) \quad \begin{cases} p_{n+1} = \operatorname{prox}_{\tau\varphi}(x_n - \tau L^* u_n) \\ x_{n+1} = P_C p_{n+1} \\ u_{n+1} = \operatorname{prox}_{\sigma\psi^*}(u_n + \sigma L(x_{n+1} + p_{n+1} - x_n)). \end{cases}$$

 $(x^k)_{k\in\mathbb{N}}\subset C$  converges to a point in  $Z\cap C=Z$ .

- In particular, we use C as the mass constraint.
- C can change deterministically/randomly among iterations.

<sup>&</sup>lt;sup>5</sup> with J. Deride, S. López Rivera, and C. Vega  $\langle \Box \rangle \langle \Box \rangle \langle \Box \rangle \langle \Xi \rangle$ 

### $\operatorname{prox}_{\gamma\varphi}$ computation

Motivation

• In all previous methods we need to compute  $\operatorname{prox}_{\gamma\varphi}(m, w)$ .

Split/Unsplit

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Numerical experiences

• Recall that  $\varphi \colon (m, w) \mapsto \sum_{i,j} \phi_{i,j}(m_{i,j}, w_{i,j})$ , where  $\phi_{i,j} \colon (\mu, \omega) \mapsto \hat{b}(\mu, \omega) + F_{i,j}(\mu)$  and  $F_{i,j} = F(x_{i,j}, \cdot)$ .

Splitting algorithms

• We have  $\operatorname{prox}_{\gamma\varphi}(m,w) = (\operatorname{prox}_{\gamma\phi_{i,j}}(m_{i,j},w_{i,j}))_{i,j}$ .

Prox computation

$$\operatorname{prox}_{\gamma\phi} \colon (\mu, \omega) \mapsto \begin{cases} (0,0), & \text{if } m \le \gamma F'(0); \\ (p^*, p^* P_K w/(p^* + \gamma)), & \text{if } m > \gamma F'(0), \end{cases}$$

where  $p^* \ge 0$  is the unique solution to

$$(p + \gamma F'(p) - m)(p + \gamma)^2 - \frac{\gamma}{2}|P_K w|^2 = 0.$$

• We extend prox in Papadakis, Peyre & Oudet (2014) used in the context of optimal transport (we include F and K).

Algorithm

Denoting  $\varphi = b_2 + F$ ,  $A = -\nu \Delta_h$ ,  $B = \operatorname{div}_h$ , the classic CP splitting reads :

$$\begin{pmatrix} u^{k+1}\\\lambda^{k+1} \end{pmatrix} = \begin{pmatrix} u^k + \gamma(A\bar{m}^k + B\bar{w}^k)\\\lambda^k + \gamma(h^2\sum_{i,j}\bar{m}^k_{i,j} - 1) \end{pmatrix}$$
$$\begin{pmatrix} m^{k+1}\\w^{k+1} \end{pmatrix} = \operatorname{prox}_{\tau\varphi} \begin{pmatrix} m^k - \tau(A^{\top}u^{k+1} + h^2\lambda^{k+1}\mathbf{1})\\w^k - \tau B^{\top}u^{k+1} \end{pmatrix}$$
$$\begin{pmatrix} \bar{m}^{k+1}\\\bar{w}^{k+1} \end{pmatrix} = \begin{pmatrix} 2m^{k+1} - m^k\\2w^{k+1} - w^k \end{pmatrix}$$

but primal iterates do not satisfy any of the constraints, which extremely affect the speed of the method.

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### Algorithm

Imposing, the constraint  $\int m = 1$  the PCP reads

$$\begin{pmatrix} u^{k+1}\\\lambda^{k+1} \end{pmatrix} = \begin{pmatrix} u^k + \gamma(A\bar{m}^k + B\bar{w}^k)\\\lambda^k + \gamma(h^2\sum_{i,j}\bar{m}^k_{i,j} - 1) \end{pmatrix}$$

$$\begin{pmatrix} n^{k+1}\\v^{k+1} \end{pmatrix} = \operatorname{prox}_{\tau\varphi} \begin{pmatrix} m^k - \tau(A^{\top}u^{k+1} + h^2\lambda^{k+1}\mathbf{1})\\w^k - \tau B^{\top}u^{k+1} \end{pmatrix}$$

$$\begin{pmatrix} m^{k+1}\\w^{k+1} \end{pmatrix} = \begin{pmatrix} \mathbf{1} + (n^{k+1} - \mathbf{1}\sum_{i,j=1}^{N_h} n^{k+1}_{i,j})\\v^{k+1} \end{pmatrix}$$

$$\begin{pmatrix} \bar{m}^{k+1}\\\bar{w}^{k+1} \end{pmatrix} = \begin{pmatrix} m^{k+1} + n^{k+1} - m^k\\w^{k+1} + v^{k+1} - w^k \end{pmatrix}$$

which is much faster, and any of the primal iterates satisfy the imposed constraint. By including splitting, we do not need any matrix inversion.

We consider the first-order stationary MFG system (Almulla-Ferreira-Gomes, 2015)

$$\frac{1}{2}|\nabla u|^2 - \lambda = \log m - \sin(2\pi x) - \sin(2\pi y),$$
$$\operatorname{div}(m\nabla u) = 0, \quad \int_{\mathbb{T}^2} m dx = 1, \quad \int_{\mathbb{T}^2} u dx = 0,$$

with explicit solution

$$u(x,y) = 0, \quad m(x,y) = e^{\sin(2\pi x) + \sin(2\pi y) - \lambda},$$
$$\lambda = \log\left(\int_{\mathbb{T}^2} e^{\sin(2\pi x) + \sin(2\pi y)} dx dy\right).$$

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47/51

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Motivation $(P_h)$ Splitting algorithmsSplit/UnsplitNumerical experiences00000000000000000000000000000000000000 $(P_h)$  $(P_h)$ </tr

Test 2: Comparison with ADMM

$$f(x, y, m) = \frac{1}{2}(m - \bar{m})$$

where  $\bar{m}$  is centered gaussian density.



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#### Also done:

- We tackle the more general case  $H \colon (x,p) \mapsto |p|^{q'}/q'$  for q > 1.
- We also present 2 additional tests.

Splitting algorithms

• Test 3: We include hard density constraints in some regions.

Split/Unsplit

Numerical experiences

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- Test 4: We vary  $q \neq 2$ .
- Dynamic case.

#### In preparation:

- More general Hamiltonians, congestion, non-local couplings.
- Other efficient projection configurations for avoiding matrix inversions.

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### Merci de votre attention !